

# SUMMARY OF RESEARCH PERFORMED FOR NATIONAL AERONAUTICS AND SPACE ADMINISTRATION

Electronics Research Division  
Communications Research Center

on the project entitled

## ARCHITECTURE AND IMPLEMENTATION CONSIDERATIONS OF A HIGH-SPEED VITERBI DECODER FOR A CHANNEL SUBCODE

Contract Number N00019-77-2938

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**Electrical Engineering Division  
Goddard Space Flight Center**

**on the project entitled**

**ARCHITECTURE AND IMPLEMENTATION  
CONSIDERATIONS OF A HIGH-SPEED VITERBI  
DECODER FOR A REED-MULLER SUBCODE**

**(Progress Report)**

**Grant Number NAG 5-2938**

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**PROGRESS REPORT OF RESEARCH PERFORMED FOR  
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Electrical Engineering Division

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**PROGRESS ON IMPLEMENTATION OF THE SUBTRELLIS IC**

Shu Lin, Gregory T. Uehara, Eric Nakamura, and Cecilia W. P. Chu

# INTRODUCTION

In this research, we have proposed the (64, 40, 8) subcode of the third-order Reed-Muller (RM) code to NASA for high-speed satellite communications. This RM subcode can be used either alone or as an inner code of a concatenated coding system with the NASA standard (255, 233, 33) Reed-Solomon (RS) code as the outer code to achieve high performance (or low bit-error rate) with reduced decoding complexity. It can also be used as a component code in a multilevel bandwidth efficient coded modulation system to achieve reliable bandwidth efficient data transmission.

This report will summarize the key progress we have made toward achieving our eventual goal of implementing a decoder system based upon this code.

In previous reports [1,6], we described results from our investigations of the complexities of various sectionalized trellis diagrams for the proposed (64, 40, 8) RM subcode. We found a specific 8-trellis diagram for this code which requires the least decoding complexity with the potential to achieve a decoding speed of 600 M bits per second (Mbps). The combination of a large number of states and a high data rate will be made possible due to the utilization of a high degree of parallelism throughout the architecture. This trellis diagram was presented and described in detail [1]. We then investigated circuit architectures to determine the feasibility of VLSI implementation of a high-speed Viterbi decoder based on this 8-section trellis diagram. We made detailed design and feasibility examinations of implementation approaches for the key blocks. Our key results for block level implementation were presented in [6].

This report will focus on our recent progress and plans regarding development of the integrated circuit prototype sub-trellis IC, particularly focusing on the design methodology.

## 1. Summary of Previous Results

We will begin this section with a brief discussion of the system block diagram in which the proposed decoder is assumed to be operating. Next, we will present some of the results from our architecture development for a sub-trellis IC which will be the basic building block for a decoder system.

### System Block Diagram

A simplified block diagram of a receiver in which the proposed decoder may be used is shown in Fig. 1. The signal enters the receiver via an antenna and is first amplified by a low noise amplifier (LNA) before being passed to the 2-PSK demodulator. We assume the functions of carrier and timing acquisition and gain control are properly performed in the demodulator. The output of the demodulator is sampled at the correct phase at the symbol rate of 960 MHz. The output of the sampler is converted to the digital domain by the 3-bit analog-to-digital converter (ADC) for decoding by the Viterbi Decoder block which follows. Our work currently focuses exclusively on the implementation of the Viterbi Decoder.

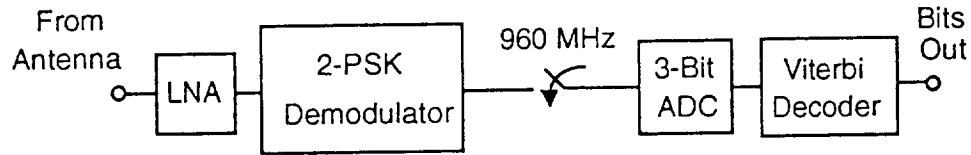
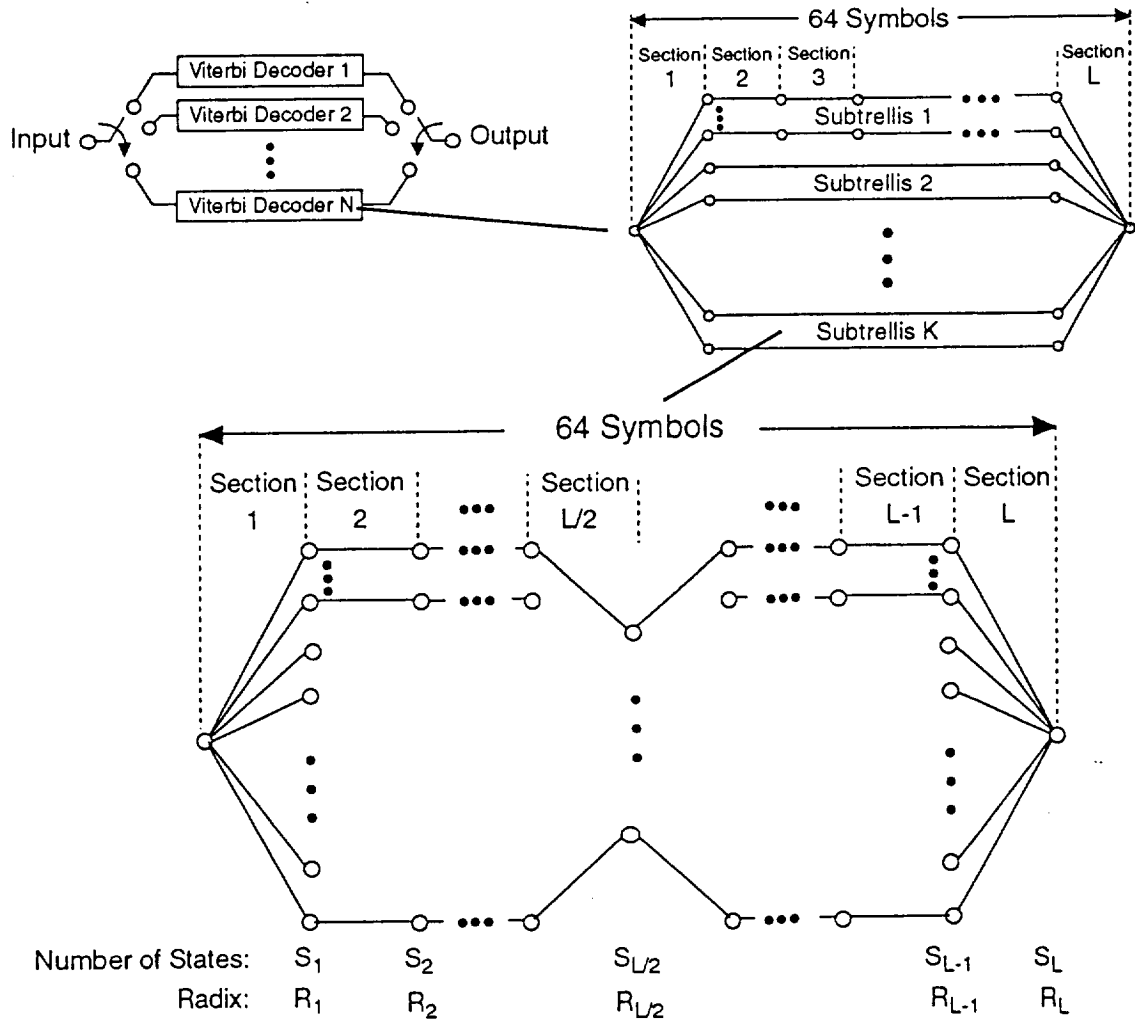


Figure 1 Block diagram of a high speed satellite receiver employing 2-PSK signalling and a Viterbi Decoder.

### Summary of System Level Architecture Design

In our earlier reports [1,6], we describe in detail the different ways in which parallelism can be utilized to decode the (64, 40) RM code. We will provide a brief summary of these descriptions in this section.

There are many diverse issues at different levels of the design requiring consideration for implementation of the (64, 40) RM code at a rate of 600 Mbits/sec. Fig. 2 illustrates the different layers of hierarchy associated with the proposed implementation. First, there are  $N$  parallel decoders with each operating on a different independent block of 64 symbols. Given a decoder which can decode a 64-symbol block at a certain rate, using  $N$  decoders and having them each operate on a different block of 64 symbols allows a throughput  $N$  times greater. Second, each decoder is implemented with  $K$  parallel isomorphic subtrellises. As described in [5], the trellis for an RM code can be decomposed into parallel isomorphic subtrellises that are connected at only the inputs and outputs as shown conceptually in Fig. 2 with  $K$  parallel subtrellises. This has a tremendous advantage for IC implementation because it minimizes the amount of routing required within the trellis which would otherwise be unrealizable at high speed for applications requiring large numbers of states. This is the key which makes an implementation using CMOS IC's at such a high rate and complexity possible. And third, there are a number of parameters associated with the implementation of each of the  $K$  subtrellises. The first is the number of sections in the subtrellis denoted as  $L$ . Next, is the number of states at the end of each section  $i$  ( $i = 1, 2, \dots, L$ ) denoted as  $|S_i|$  which will generally not be the same. Finally, there is the radix of each section denoted as  $R_i$  for radix  $R$  in section  $i$ . As the number of sections  $L$  decreases, the complexity of each section and the number of parallel branches per section increases. These trade-offs are discussed in detail in [1].



**Figure 2** Levels of hierarchy in the proposed Viterbi decoder implementation. (a) Parallel Viterbi decoders operating on different blocks of data. (b) Implementation with K parallel isomorphic subtrellises. (c) Subtrellis implementation.

After examining a number of various permutations of N, K, L, S, and R, we settled on a solution with the detailed structure shown in Fig. 3. We call this structure Trellis 2 and each Viterbi Decoder of Fig. 2b will have K equal 32. In this solution, our design goal is to meet the speed objectives in a currently available CMOS technology with N equal 2.

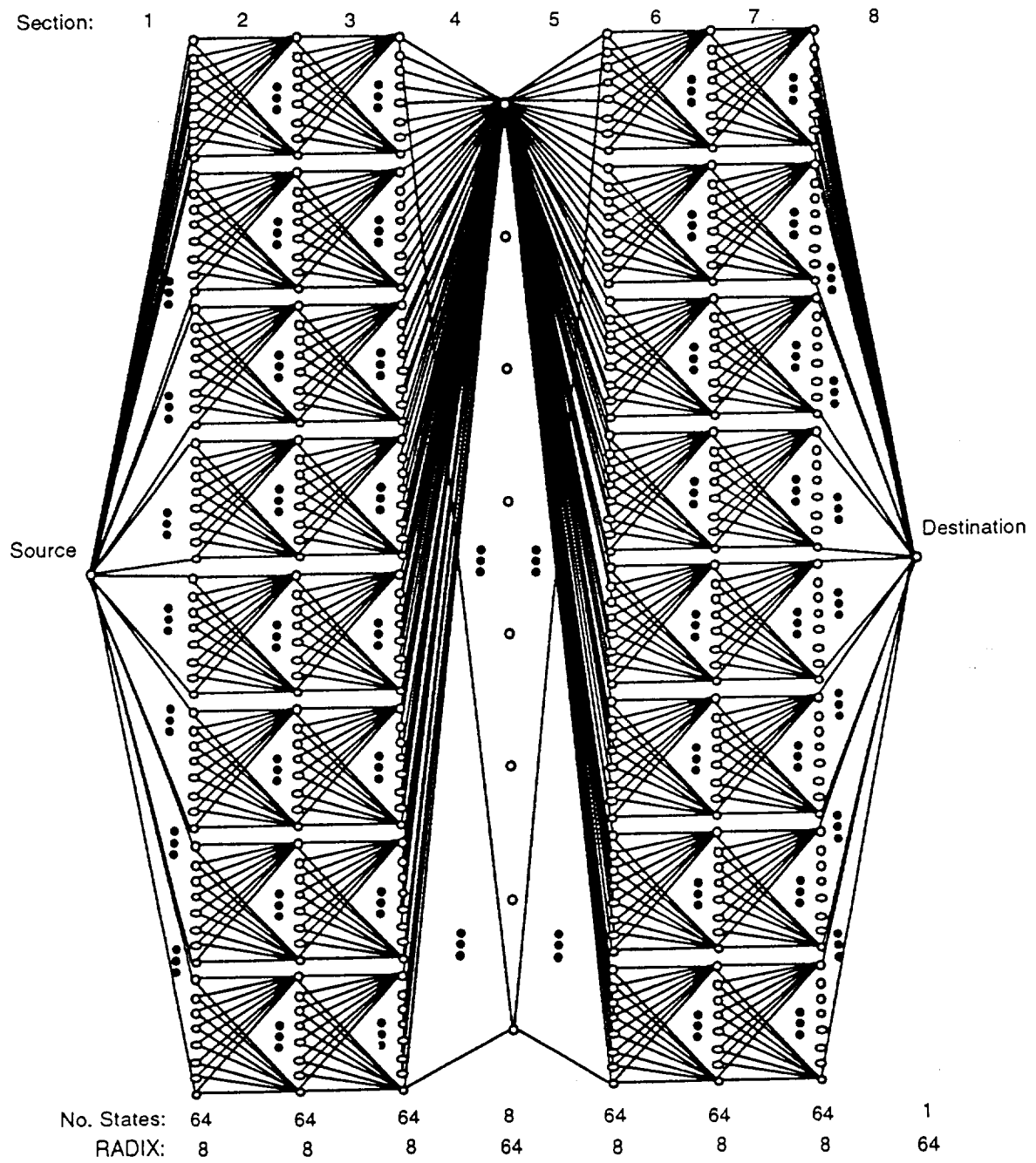


Figure 3 Detailed subtrellis structure for Trellis 2.

The key to the implementation of a (64, 40) RM decoder will be the successful implementation of an IC implementing the subtrellis shown in Fig. 3.

The key objectives of the subtrellis IC implementation are to:

1. Maximize the efficiency as measured by maximizing the utilization of the hardware (in other words, attempt to minimize the time the majority of the hardware is not being used).
2. Use a chip plan which minimizes the area used for routing (routing area is simply an overhead which should be minimized).
3. Approach the speed of 600 Mbits/sec with 2 parallel decoders.
4. Consider reliability and robustness issues. In particular, use the lowest speed system clock possible which allows high speed operation in order to reduce the number of issues which can limit the performance (which in this case would be clock skew between chips or race conditions both within and between the different ICs).
5. Consider the board design and the numbers of inputs and outputs to each chip to facilitate implementation of the final decoder system.
6. Keep the size of the IC on the order of 10 mm per side to facilitate its implementation and yield for testing.

## 2. Recent Results

Our recent efforts have focused on the design and development of the prototype IC. The goal of this portion of the project is to design and layout the circuits in a computer aided environment to create a database which can be used by a fabrication facility to generate the necessary masks to fabricate the prototype IC.

### Initial Design Procedure

The design procedure for development of the prototype IC we have used to date is as follows:

1. Block Level Design -- Define the functional performance the major blocks.
2. Timing Diagram Design -- Define the flow of data through the chip based upon the design of the major blocks.
3. Circuit Design -- Design circuits first at the gate level and then the transistor level depending upon the particular transistor logic style (complementary, pass-gate, dynamic, etc...) to perform the desired functions and at the desired speed.
4. Circuit Layout -- Create a full custom layout by hand which defines the location and size of the geometries which become the mask set for fabricating the chip.
5. Verification -- Verify that the layout performs the functions desired by the circuit through extracting the connectivity and transistor geometry information and using the file as part of an input control file for a circuit simulator, namely SPICE.
6. Full-Chip Layout -- Repeat Steps 3 - 5 outlined above for each of the cells which make up the sub-blocks which are then connected to make up the major blocks. Following this, perform these steps again as the major blocks are connected and then verified. Finally, verify the functionality of the entire layout.
7. Send a layout file for fabrication to an appropriate facility.



### Chip Plan -- Block Level Overview

An outline of the overall block plan is shown in Fig. 4. The *Clock Generation and Control* block will generate the necessary clock phases to clock the chip. Input data will enter the *Branch Metric Unit (BMU)* which will generate the branch metrics for the *Add-Compare-Select Unit (ACSU)*. The outputs of the ACSU include the winning path metrics and the winning branch labels. These are input to the *Decoder* Unit which determines the most likely path through the subtrellis for the 64-symbol block. Pipelining is used extensively within the BMU, ACSU, and the Decoder. Due to the use of block processing, we currently plan to have the input clock to the chip clock at a 60 MHz rate. The design goal is to have each IC process data at a 480 Msymbol/sec rate (300 Mbps).

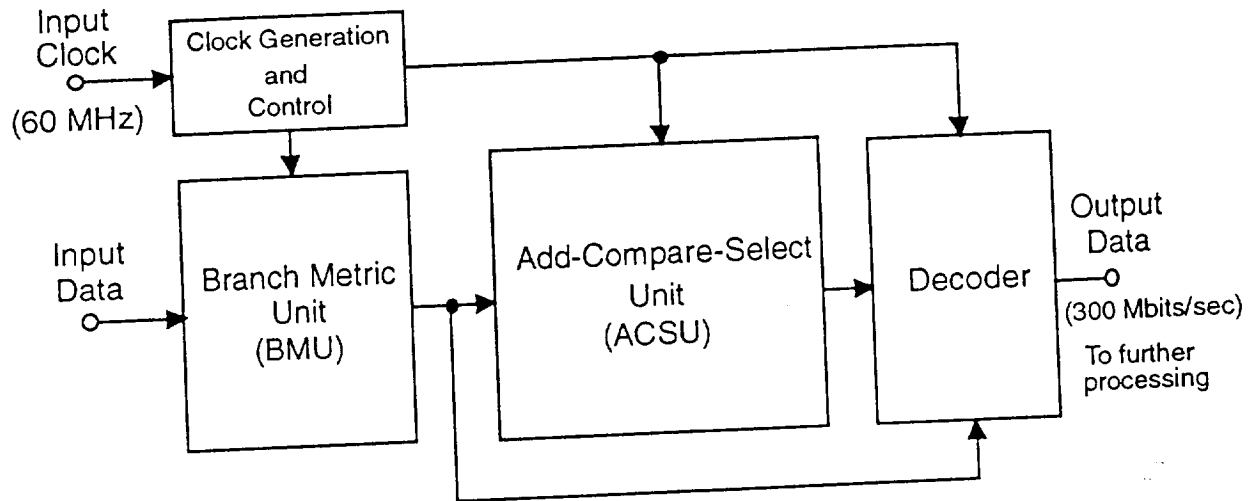


Figure 4 (a) Block diagram of the IC being developed to implement a subtrellis.

Using the Initial Design Procedure outlined above, we came up with an initial functional design for the entire chip. We completed Steps 1 - 3 for the major blocks shown in Fig. 4. We then developed a layout for an 8-Way ACS Cell which will be the basic building block in the ACSU and other key building blocks which will be repeated due to the modularity of the design. This allowed us to develop the estimates for the chip layout of a subtrellis IC shown in Table 1. Layouts and simulations assumed the use of a 0.6  $\mu\text{m}$  double-metal CMOS technology. With pads and routing, a conservative estimate of the die size in this technology is 1.2 mm x 1.2 mm.

Table 1. Estimates Using Initial Design Procedure

Block	Transistor Count	Block Size ( $\text{mm}^2$ )
Clock Generation and Control	1,000	500 $\mu\text{m}$ x 1000 $\mu\text{m}$
Branch Metric Unit	2,500	1,000 $\mu\text{m}$ x 1,000 $\mu\text{m}$
Add-Compare-Select Unit	275,000	8,000 $\mu\text{m}$ x 11,000 $\mu\text{m}$
Decoder	175,000	2,500 $\mu\text{m}$ x 11,000 $\mu\text{m}$

### Limitations of the Initial Design Procedure

The procedure outlined above was very useful for us to obtain estimates of the size and complexity of each of the blocks. However, it underscored a basic limitation of our Initial Design Procedure which is the adequate full-chip verification of the layout of a chip with nearly 500,000 transistors. This procedure was used successfully in other university projects in high speed decoders [2 - 4]. However, the overall chip complexity in our project has turned out to be significantly greater than that of this other previous work.

### LSI Logic (San Jose, California) Association, Relationship, and Support

LSI Logic is a company based in San Jose, California focused on integrated circuit development for high performance communication systems. One of the primary products of LSI Logic is a design methodology which allows customers to design and develop custom integrated circuits in a systematic and proven manner. Using this LSI Logic design methodology, customers begin with a set of functional specifications and end with prototype IC devices in near state-of-the-art CMOS technologies.

This summer, our two students involved with the development of the prototype IC, Eric Nakamura and Cecilia Chu, are spending the summer at LSI Logic as Temporary Employees. As a result of the combination of coding and VLSI development research work here at the University of Hawaii, LSI Logic has started what is planned to be a long term support relationship of our University. In this relationship, LSI Logic will supply their design methodology and chip fabrication services to the University of Hawaii in return for research updates (as is available to all companies). While there are other benefits to LSI Logic such as the potential for student hires through internship experiences, research updates will probably happen in a manner more timely than might be considered typical due to the established relationship with faculty and students. In the longer term, LSI Logic plans to support increasing amounts of research in coding and VLSI development here at the University.

Our students this summer are focusing their efforts on learning the LSI Logic methodology in the context of a development project for LSI Logic as they are hired as Temporary Employees. In the coming fall, they will bring back to the University the LSI Logic methodology which we plan to use for development of the prototype subtrellis IC. The advantages for our current project are three-fold.

First, LSI Logic has currently available a 0.35  $\mu\text{m}$  CMOS technology which would result in a nearly a 4 times reduction in the layout size of a given cell. As a result, we would probably tend toward the use of standard cells from the LSI Logic family as opposed to a full custom hand design as we had planned. This takes us to the second point. In the LSI Logic design methodology, functional blocks are described using a Hardware Description Language (HDL). This programming language is first used to describe a given function and which can be simulated and verified. This HDL can then be used to directly develop a circuit layout which would be a connection of standard cells from a cell library which would implement the described function. This will greatly reduce development time as compared with our Initial Design Procedure. While the chip area is greater when standard cells are used to implement a given function as compared with a full custom hand design<sup>1</sup>, the use of a more aggressive technology can still result in a decrease in the overall chip size. And third and most important, the LSI Logic design methodology has been used effectively to develop integrated circuits with more than one million transistors. This proven method will greatly increase the probability of working devices in our initial design. It will allow us to verify the circuit layout with a much greater confidence as compared with the Initial Design Procedure.

We expect that the design results from the Initial Design Procedure, which provided an initial design of the prototype IC through to the circuit level, will be extremely beneficial to our development using the

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1. *Hand design* signifies a layout drawn by hand on a computer as opposed to one automatically generated on a computer.

LSI Logic design methodology. We are aware of the critical blocks and now have points of reference from both the circuit and layout standpoints with which to compare our new designs. This should result in a superior solution than would otherwise have been obtained if either of the design approaches were used exclusively.

### 3. Summary

In our recent efforts, we completed the majority of the circuit design down to the circuit level using what we call our Initial Design Procedure. Through this process, we believe a prototype IC which can be used to implement the 600 Mbps decoder is achievable using a 0.6  $\mu\text{m}$  CMOS technology using the approaches described in our previous report. This summer we have our two students Eric Nakamura and Cecilia Chu at LSI Logic, learning the LSI Logic Design Methodology which they plan to bring back to the University. We believe the LSI Logic design methodology will result in a circuit layout whose performance can be better verified prior to fabrication than would otherwise be possible using our Initial Design Procedure. Thus, the initial prototype circuits will have a much greater chance of functionality. LSI Logic intends to be involved with the fabrication of the prototype IC using their 0.35  $\mu\text{m}$  CMOS technology. This newly developed association between LSI Logic and the University of Hawaii should prove to be very useful as we progress in our development of the prototype IC on our way to building a prototype decoder system.

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